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MySQL 8.0:

What Is New in Optimizer and Executor?

Norvald H. Ryeng
Software Development Senior Manager
MySQL Optimizer Team
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Program Agenda

- ▶ Common table expressions
- ▶ Window functions
- ▶ UTF-8 support
- ▶ GIS
- ▶ SKIP LOCKED, NOWAIT
- ▶ JSON functions
- ▶ Index extensions
- ▶ Cost model
- ▶ Hints

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CTEs As Alternative to Derived Tables

- A derived table is a subquery in the FROM clause:

```
SELECT ... FROM (subquery) AS derived, t1 ...
```

- A CTE is just like a derived table, but the declaration comes before the query block:

```
WITH derived AS (subquery)  
SELECT ... FROM derived, t1 ...
```

- May precede SELECT, UPDATE, DELETE and be used in subqueries:

```
WITH derived AS (subquery)  
DELETE FROM t1 WHERE t1.a IN (SELECT b FROM derived);
```

CTEs vs. Derived Tables

- Better readability
- Can be referenced multiple times
- Can refer to other CTEs
- Improved performance

CTEs Provide Better Readability

Derived table:

```
SELECT ...  
FROM t1 LEFT JOIN ((SELECT ... FROM ...)) AS dt JOIN t2 ON ...) ON ...
```

CTE:

```
WITH dt AS (SELECT ... FROM ...)  
SELECT ...  
FROM t1 LEFT JOIN (dt JOIN t2 ON ...) ON ...
```


CTEs Can Be Referenced Multiple Times

- Derived tables can't be referenced twice

```
SELECT ...  
FROM (SELECT a, b, SUM(c) AS s FROM t1 GROUP BY a, b) AS d1  
JOIN (SELECT a, b, SUM(c) AS s FROM t1 GROUP BY a, b) AS d2 ON d1.b = d2.a;
```

- CTEs can

```
WITH d AS (SELECT a, b, SUM(c) AS s FROM t1 GROUP BY a, b)  
SELECT ... FROM d AS d1 JOIN d AS d2 ON d1.b = d2.a;
```

- Better performance with materialization
 - Multiple CTE references are only materialized once
 - Derived tables and views will be materialized once per reference

CTEs Can Refer to Other CTEs

- Derived tables can't refer to other derived tables

```
SELECT ...  
FROM (SELECT ... FROM ...) AS d1, (SELECT ... FROM d1 ...) AS d2 ...  
ERROR: 1146 (42S02): Table 'db.d1' doesn't exist
```

- CTEs can refer other CTEs

```
WITH d1 AS (SELECT ... FROM ...),  
     d2 AS (SELECT ... FROM d1 ...)  
SELECT ...  
FROM d1, d2 ...
```

Recursive CTEs

WITH RECURSIVE *cte* AS

(**SELECT ... FROM *table_name***

UNION [DISTINCT|ALL]

SELECT ... FROM *cte*, *table_name*)

SELECT ... FROM *cte*;

/* "seed" SELECT */

/* "recursive" SELECT */

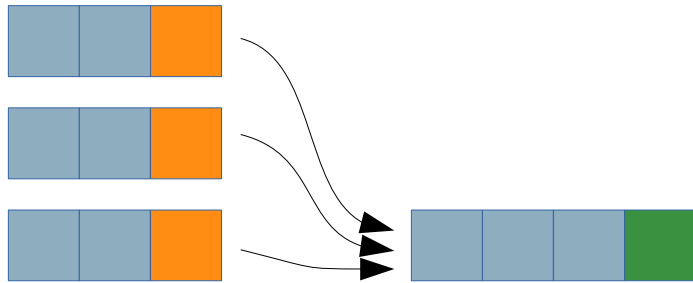
- A recursive CTE refers to itself in a subquery
- The "seed" SELECT is executed once to create the initial data subset
- The "recursive" CTE is executed repeatedly
 - Stops when iteration doesn't generate any new rows
 - Set `cte_max_recursion_depth` to limit recursion
- Useful to traverse hierarchies (parent/child, part/subpart)

Program Agenda

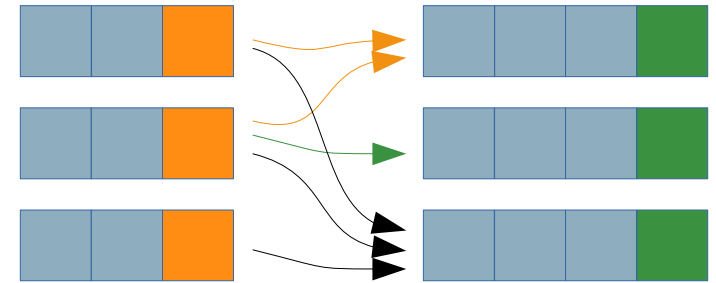
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Window Functions

- Similar to aggregation functions
 - Computes one value based on multiple rows
 - But does not group



Aggregation function



Window function

Window Function Example

Sum up total salary per department

```
SELECT name, dept_id, salary,  
       SUM(salary) OVER (PARTITION BY  
       dept_id) AS dept_total  
FROM employee  
ORDER BY dept_id, name;
```

The **OVER**
keyword signals a
window function

PARTITION ⇒
disjoint set of
rows in result set

name	dept_id	salary
Newt	NULL	75000
Dag	10	NULL
Ed	10	100000
Fred	10	60000
Jon	10	60000
Michael	10	70000
Newt	10	80000
Lebedev	20	65000
Pete	20	65000
Jeff	30	300000
Will	30	70000

Window Function Example

Sum up total salary per department

```
SELECT name, dept_id, salary,  
       SUM(salary) OVER (PARTITION BY  
                          dept_id) AS dept_total  
FROM employee  
ORDER BY dept_id, name;
```

The **OVER**
keyword signals a
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name	dept_id	salary	dept_total
Newt	NULL	75000	75000
Dag	10	NULL	370000
Ed	10	100000	370000
Fred	10	60000	370000
Jon	10	60000	370000
Michael	10	70000	370000
Newt	10	80000	370000
Lebedev	20	65000	130000
Pete	20	65000	130000
Jeff	30	300000	370000
Will	30	70000	370000

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UTF-8 Support

- Collations based on Unicode 9.0
- utf8mb4 as default character set
 - utf8mb4_0900_ai_ci as default collation
- Accent and case sensitive collations
 - Including 20+ language specific collations
 - Now also Japanese and Russian
- Significantly improved performance 👍
- Unicode support in regular expressions



New Default Character Set

- No change to existing tables
- Only has effect on new tables/schemas when character set is not explicitly specified

Upgrade recommendations

- Upgrade first, change character set and collation afterwards
- Don't mix collations
 - Error: "Illegal mix of collations"
 - Slower queries because indexes can't be used

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GIS

- Geography support
 - Latitude-longitude on an ellipsoid
- Spatial reference systems (SRSs)
 - More than 5000 SRSs from the EPSG Geodetic Parameter Dataset
 - INFORMATION_SCHEMA.ST_SPATIAL_REFERENCE_SYSTEMS
 - CREATE/DROP SPATIAL REFERENCE SYSTEM
- SRID aware spatial datatypes, indexes and functions
- ST_Transform to convert between SRSs — **New in MySQL 8.0.13!**
- Functions and expressions as default values — **New in MySQL 8.0.13!**
 - Including ST_GeomFromText

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SELECT ... FOR UPDATE SKIP LOCKED

Problem: Hot row contention, multiple worker threads accessing the same rows.

SKIP LOCKED Solution:

- Only read rows that aren't locked
- InnoDB skips a locked row, and the next one goes to the result set

Example:

- Booking system: Skip orders that are pending

```
START TRANSACTION;
```

```
SELECT * FROM seats WHERE seat_no BETWEEN 2 AND 3 AND booked = 'NO'
```

```
FOR UPDATE SKIP LOCKED;
```

SELECT ... FOR UPDATE NOWAIT

Problem: Hot row contention, multiple worker threads accessing the same rows.

NOWAIT Solution:

- Fail immediately if any rows are locked
- Without NOWAIT, the query waits for `innodb_lock_wait_timeout` (default 50 seconds) before giving up

Example:

- Booking system: Fail if rows are locked

```
START TRANSACTION;  
SELECT * FROM seats WHERE seat_no BETWEEN 2 AND 3 AND booked = 'NO'  
FOR UPDATE NOWAIT;  
ERROR 3572 (HY000): Statement aborted because lock(s) could not be acquired ...
```

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JSON Functions

JSON_ARRAY_APPEND()

JSON_ARRAY_INSERT()

JSON_ARRAY()

JSON_CONTAINS_PATH()

JSON_CONTAINS()

JSON_DEPTH()

JSON_EXTRACT()

JSON_INSERT()

JSON_KEYS()

JSON_LENGTH()

JSON_MERGE[_PRESERVE]()

JSON_OBJECT()

JSON_QUOTE()

JSON_REMOVE()

JSON_REPLACE()

JSON_SEARCH()

JSON_SET()

JSON_TYPE()

JSON_UNQUOTE()

JSON_VALID()

JSON_PRETTY()

JSON_STORAGE_SIZE()

JSON_STORAGE_FREE()

JSON_ARRAYAGG()

JSON_OBJECTAGG()

JSON_MERGE_PATCH()

JSON_TABLE()

JSON Data Type

```
CREATE TABLE t1 (json_col JSON);
```

```
INSERT INTO t1 VALUES (  
  '{ "people": [  
    { "name": "John Smith", "address": "780 Mission St, San Francisco, CA 94103"},  
    { "name": "Sally Brown", "address": "75 37th Ave S, St Cloud, MN 94103"},  
    { "name": "John Johnson", "address": "1262 Roosevelt Trail, Raymond, ME 04071"}  
  ] }'  
);
```

JSON_TABLE

Convert JSON documents to relational tables

```
SELECT people.*
FROM t1,
     JSON_TABLE(json_col, '$.people[*]' COLUMNS (
        name VARCHAR(40) PATH '$.name',
        address VARCHAR(100) PATH '$.address')) AS people;
```

name	address
John Smith	780 Mission St, San Francisco, CA 94103
Sally Brown	75 37th Ave S, St Cloud, MN 9410
John Johnson	1262 Roosevelt Trail, Raymond, ME 04071

JSON_TABLE

Filter JSON data on SQL expression

```
SELECT people.*
FROM t1,
     JSON_TABLE(json_col, '$.people[*]' COLUMNS (
         name VARCHAR(40) PATH '$.name',
         address VARCHAR(100) PATH '$.address')) AS people
WHERE people.name LIKE 'John%';
```

name	address
John Smith	780 Mission St, San Francisco, CA 94103
John Johnson	1262 Roosevelt Trail, Raymond, ME 04071

JSON_TABLE Nested Arrays

```
[
  { "father": "John", "mother": "Mary",
    "marriage_date": "2003-12-05",
    "children": [ { "name": "Eric", "age": 12 },
                  { "name": "Beth", "age": 10 } ]
  },
  { "father": "Paul", "mother": "Laura",
    "children": [ { "name": "Sarah", "age": 9 },
                  { "name": "Noah", "age": 3 },
                  { "name": "Peter", "age": 1 } ]
  }
]
```

id	father	married	child_id	child	age
1	John	1	1	Eric	12
1	John	1	2	Beth	10
2	Paul	0	1	Sarah	9
2	Paul	0	2	Noah	3
2	Paul	0	3	Peter	1

JSON_TABLE Nested Arrays

```
JSON_TABLE (families, '$[*]' COLUMNS (  
  id FOR ORDINALITY,  
  father VARCHAR(30) PATH '$.father',  
  married INTEGER EXISTS PATH  
    '$.marriage_date',  
  NESTED PATH '$.children[*]' COLUMNS (  
    child_id FOR ORDINALITY,  
    child VARCHAR(30) PATH '$.name',  
    age INTEGER PATH '$.age'  
  )  
))
```

id	father	married	child_id	child	age
1	John	1	1	Eric	12
1	John	1	2	Beth	10
2	Paul	0	1	Sarah	9
2	Paul	0	2	Noah	3
2	Paul	0	3	Peter	1

JSON Array Aggregation

```
CREATE TABLE t1 (id INT, grp INT,  
jsoncol JSON);
```

```
INSERT INTO t1 VALUES (1, 1,  
'{"key1":"value1","key2":"value2"}');
```

```
INSERT INTO t1 VALUES (2, 1,  
'{"keyA":"valueA","keyB":"valueB"}');
```

```
INSERT INTO t1 VALUES (3, 2,  
'{"keyX":"valueX","keyY":"valueY"}');
```

```
SELECT JSON_ARRAYAGG(jsoncol)  
FROM t1;
```

```
[{"key1":"value1","key2":"value2"},  
{"keyA":"valueA","keyB":"valueB"},  
{"keyX":"valueX","keyY":"valueY"}]
```

Yet to come: JSON_ARRAYAGG as window function

JSON Object Aggregation

```
CREATE TABLE t1 (id INT, grp INT,  
jsoncol JSON);
```

```
INSERT INTO t1 VALUES (1, 1,  
'{"key1":"value1","key2":"value2"}');
```

```
INSERT INTO t1 VALUES (2, 1,  
'{"keyA":"valueA","keyB":"valueB"}');
```

```
INSERT INTO t1 VALUES (3, 2,  
'{"keyX":"valueX","keyY":"valueY"}');
```

```
SELECT grp, JSON_OBJECTAGG(id, jsoncol)  
FROM t1  
GROUP BY grp;
```

1	<pre>{"1":{"key1":"value1","key2":"value2"}, "2":{"keyA":"valueA","keyB":"valueB"}}</pre>
2	<pre>{"3":{"keyX":"valueX","keyY":"valueY"}}</pre>

Yet to come: JSON_OBJECTAGG as window function

Partial Update of JSON

- Automatically detect potential in-place updates
 - JSON_SET
 - JSON_REPLACE
 - JSON_REMOVE
 - Source and target columns are the same
 - Updating existing value, not adding new one
 - Sufficient space is available
 - Smaller than or same size as some old value
- Only the diff is sent in binlog

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Invisible Index

- Index is maintained by the storage engine, but invisible to the optimizer
 - Override with SET optimizer_switch='use_invisible_indexes=on';
- Primary key can't be invisible
- Use case: Check for performance drop *before* removing index

```
ALTER TABLE t1 ALTER INDEX idx INVISIBLE;  
SHOW INDEXES FROM t1;
```

Table	Key_name	Column_name	Visible
t1	idx	a	NO

Descending Index

```
CREATE TABLE t1 (  
  a INT,  
  b INT,  
  INDEX a_b (a DESC, b ASC)  
);
```

- In 5.7: Index in ascending order is created, server scans it backwards
- In 8.0: Index in descending order is created, server scans it forwards
 - Works on B-tree indexes only
 - Benefits:
 - Use indexes instead of filesort for ORDER BY clause with ASC/DESC sort key
 - Forward index scan is slightly faster than backward index scan

Functional Index — **New in 8.0.13**

- Index over an expression

```
CREATE TABLE t1 (col1 INT, col2 INT, INDEX func_index ((ABS(col1))));
```

```
CREATE INDEX idx1 ON t1 ((col1 + col2));
```

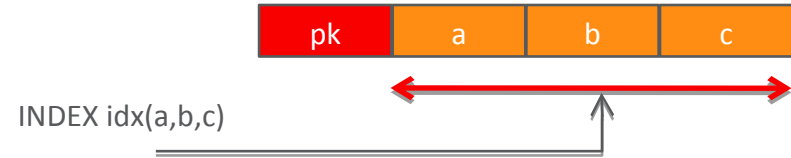
```
CREATE INDEX idx2 ON t1 ((col1 + col2), (col1 - col2), col1);
```

```
ALTER TABLE t1 ADD INDEX ((col1 * 40));
```

- No longer necessary to create an indexed generated column
- Not allowed as PRIMARY KEY

Index Skip-Scan — **New in 8.0.13**

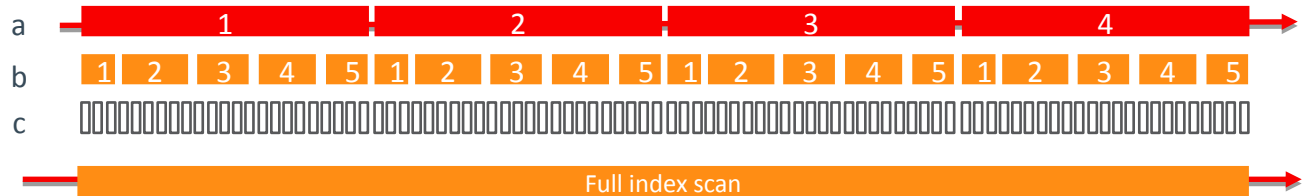
- Used for multi-column index
- Condition not on first part of index
- Cost-based choice



SELECT * FROM t1 WHERE b < 3;

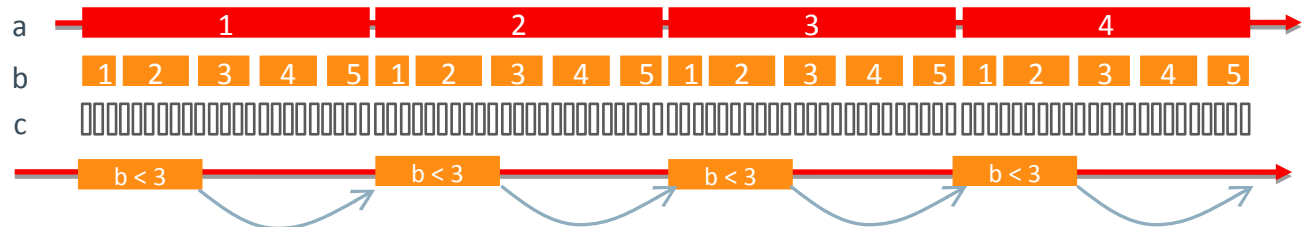
Before MySQL 8.0.13:

Full index scan



MySQL 8.0.13:

Skip-scan



**Thanks for the patch,
Facebook!**



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Motivation for Improving the MySQL Cost Model

- Produce more correct cost estimates
 - Better decisions by the optimizer should improve performance
- Adapt to new hardware architectures
 - SSD, larger memories, caches
- More maintainable cost model implementation
 - Avoid hard coded “cost constants”
 - Refactoring of existing cost model code
- Configurable and tunable
- Make more of the optimizer cost-based



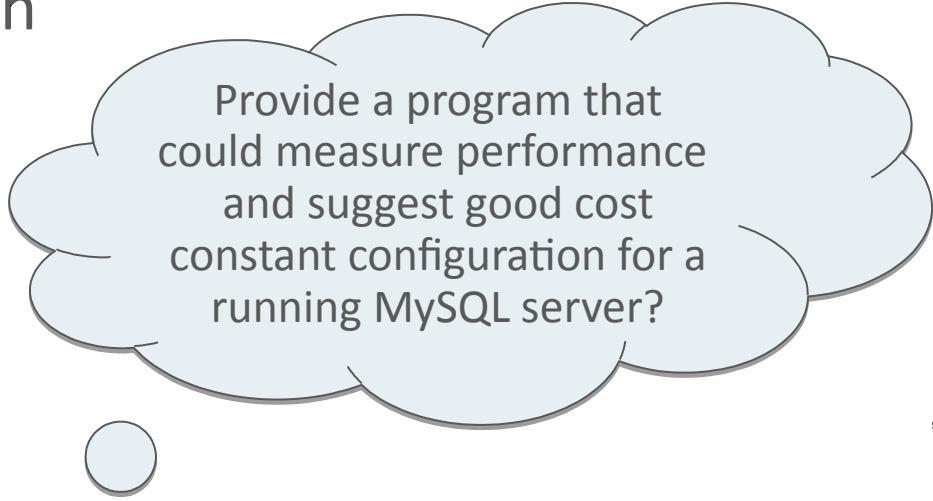
**Faster
queries**

New Storage Technologies

- Time to do a table scan of 10 million records:

Memory	5 s
SSD	20 - 146 s
Hard disk	32 - 1465 s

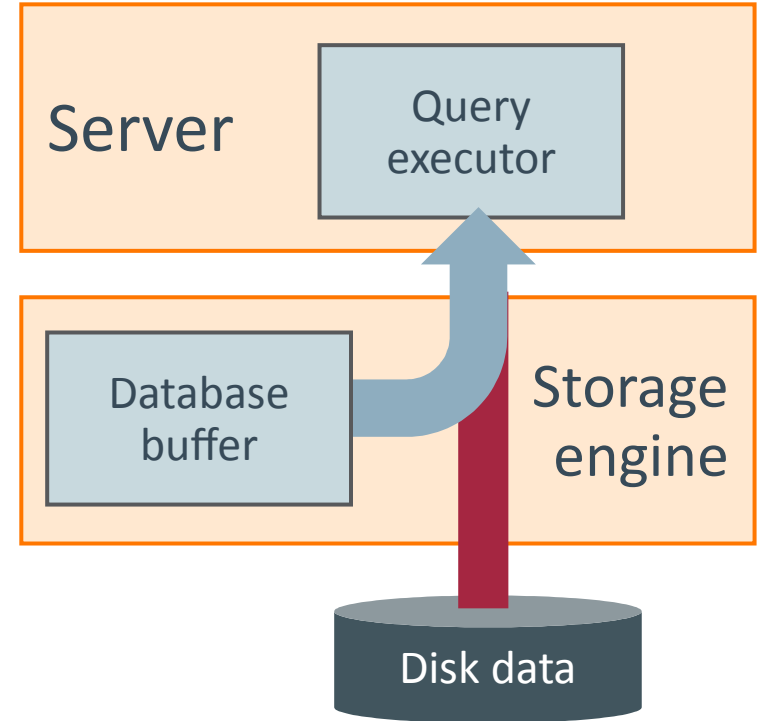
- Adjust cost model to support different storage technologies
- Provide configurable cost constants for different storage technologies



Provide a program that could measure performance and suggest good cost constant configuration for a running MySQL server?

Memory Buffer Aware Cost Estimates

- Storage engines:
 - Estimate for how much of data and indexes are in a memory buffer
 - Estimate for hit rate for memory buffer
- Optimizer cost model:
 - Take into account whether data is already in memory or need to be read from disk



Histograms

- Provides the optimizer with information about column value distribution
- To create/recalculate histogram for a column:
`ANALYZE TABLE table UPDATE HISTOGRAM ON column WITH n BUCKETS;`
- May use sampling
 - Sample size is based on available memory (histogram_generation_max_mem_size)
- Automatically chooses between two histogram types:
 - Singleton: One value per bucket
 - Equi-height: Multiple values per bucket

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Join Order Hints

- No need to reorganize the FROM clause to add join order hints like you will for STRAIGHT_JOIN
- `/*+ JOIN_ORDER(t1, t2, ...) */`
- `/*+ JOIN_PREFIX(t1, t2, ...) */`
- `/*+ JOIN_SUFFIX(..., ty, tz) */`
- `/*+ JOIN_FIXED_ORDER() */`
 - Replacement for STRAIGHT_JOIN

View/Derived Table/CTE Merge Hints

- Views, derived tables and CTEs are, if possible, merged into outer query
- NO_MERGE can override default behavior

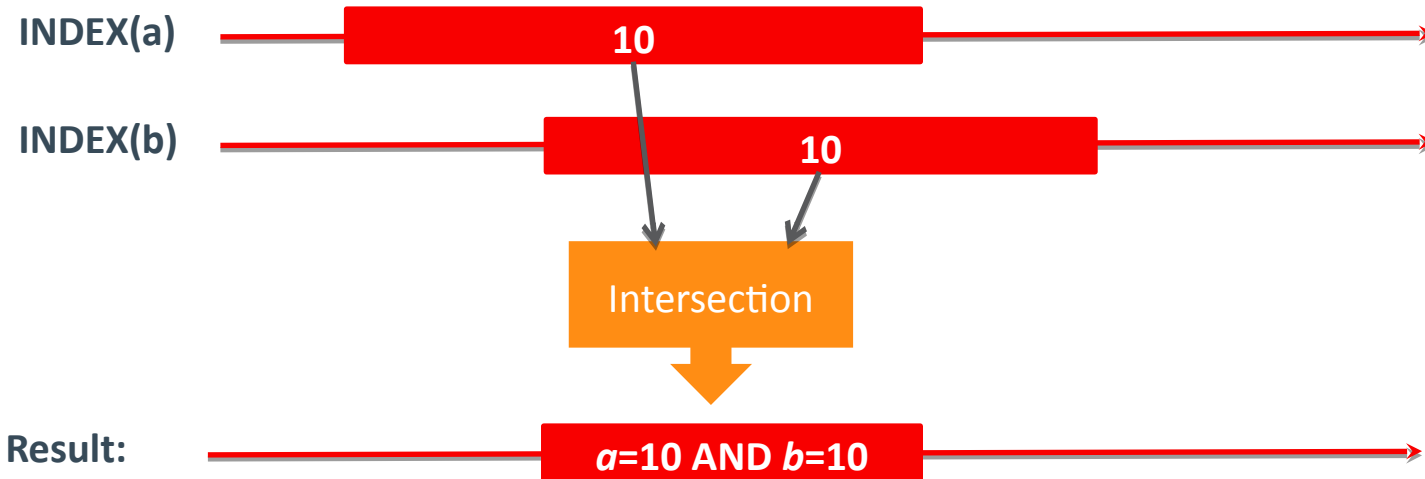
```
SELECT /*+ NO_MERGE(dt) */ *  
FROM t1 JOIN (SELECT x, y FROM t2) dt ON t1.x = dt.x;
```

- MERGE will force a merge

```
SELECT /*+ MERGE(dt) */ *  
FROM t1 JOIN (SELECT x, y FROM t2) dt ON t1.x = dt.x;
```

Index Merge Execution

SELECT * FROM *t1* WHERE *a*=10 AND *b*=10



Index Merge Hints

- Index merge: Merge rows from multiple range scans on a single table
- Algorithms: union, intersection, sort union
- Users can specify which indexes to use for index merge
 - `/*+ INDEX_MERGE(t1 idx1, idx2, ...) */`
 - `/*+ NO_INDEX_MERGE(t1 idx1, idx2, ...) */`

Hints to Set Session Variables

- Set a session variable for the duration of a single statement
- Examples:

```
SELECT /*+ SET_VAR(sort_buffer_size = 16M) */ name FROM people ORDER BY name;
```

```
INSERT /*+ SET_VAR(foreign_key_checks = OFF) */ INTO t2 VALUES (1, 1), (2, 2), (3, 3);
```

```
SELECT /*+ SET_VAR(optimizer_switch = 'condition_fanout_filter = off') */ *  
FROM customer JOIN orders ON c_custkey = o_custkey  
WHERE c_acctbal < 0 AND o_orderdate < '1993-01-01';
```

- Note: Not all session variables are settable through hints:

```
mysql> SELECT /*+ SET_VAR(max_allowed_packet=128M) */ * FROM t1;
```

```
Empty set, 1 warning (0,00 sec)
```

```
Warning (Code 4537): Variable 'max_allowed_packet' cannot be set using SET_VAR hint.
```

Feature descriptions and design details
directly from the source.

<http://mysqlserverteam.com/>

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